

# L-Store: Milestone 1

ECS 165A: Database Systems

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# 3 Main Parts



**Data Model**

Columnar Data Storage

Base Page vs Tail Page

Page Range



**Bufferpool  
Management**

Page Directory

Index Directory



**Query Interface**

INSERT

UPDATE

SELECT

DELETE & SUM

# Data Model

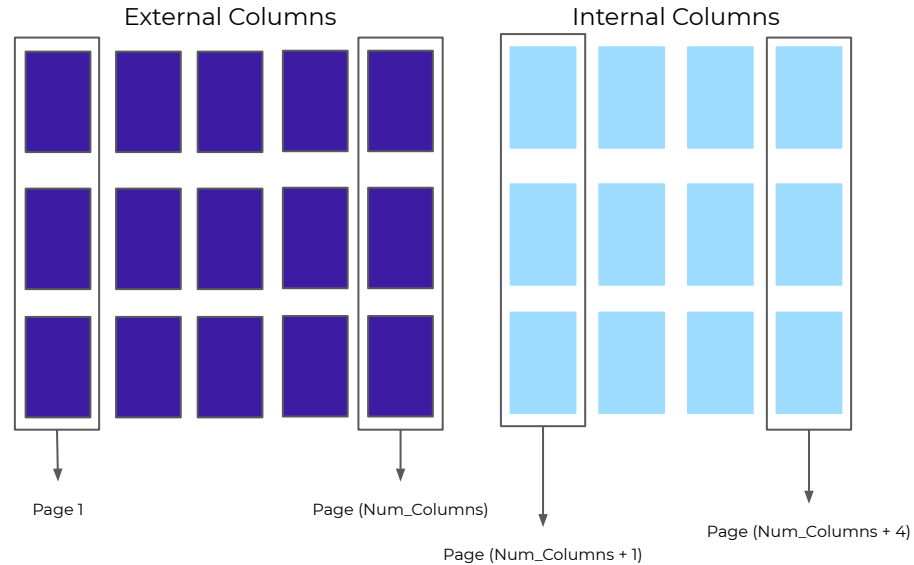
# Columnar Data Storage



External = the records



Internal = meta records  
(RID, SE, IND, TIME)



$$\text{Num\_Total\_Columns} = \text{Num\_Columns} + 4$$



# Base Page vs Tail Page



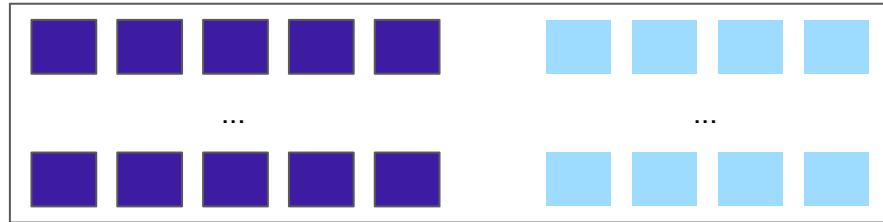
External = the records



Internal = meta records  
(RID, SE, IND, TIME)

External Columns

Internal Columns



→ **Base Page**

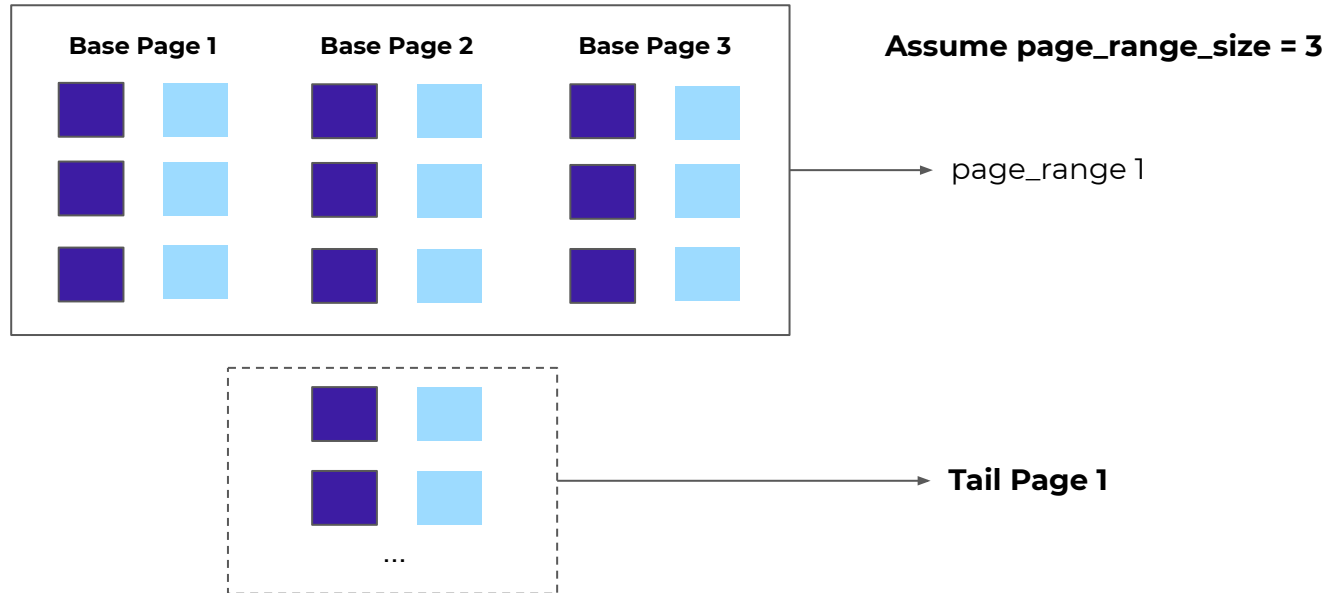


→ **Tail Page**  
depends on update

# Page Range Implementation



External = the records      Internal = meta records (RID, SE, IND, TIME)



# Helper Functions: Page

- In `page.py`:
  - `has_capacity()`: This function make sure the page still has capacity to add with increment of 8 (bytes).
  
- In `table.py`:
  - `checker()`: This function uses `has_capacity()` in order to add pages when capacity of particular page is full.

# Helper Functions: Read, Write and Edit

```
def read(offset)
```

```
    # use in get_schema_encoding and get_indirection
```

```
    return data[offset*8: (offset+1)*8]
```

```
def write(offset, value)
```

```
    # use to write into both base and tail pages
```

```
    data[self.num_records * 8: (self.num_records + 1) * 8] =  
    value.to_bytes(8, byteorder='big')  
    num_records += 1
```

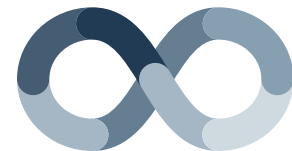
```
def edit(offset, value)
```

```
    # use in set_schema_encoding and set_indirection
```

```
    data[offset*8: (offset+1)*8] = value.to_bytes(8, byteorder = 'big')
```



# Bufferpool Management



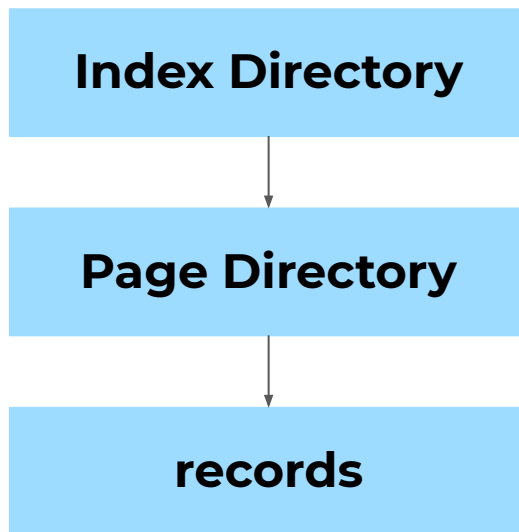
# Page vs Index Directory

	<b>Page Directory</b>	<b>Index Directory</b>
<b>Data Structure</b>	Hashmap (Dict)	Hashmap (Dict)
<b>Key</b>	RID	KEY
<b>Value</b>	PageID, Offsets	RID



# Page vs Index Directory

In **Milestone 1**, we did not use `column` parameter since we assume that the key is at column 0



# Query Interface

# Helper Functions: Get and Set + Exist

## Getters:

- `get_schema_encoding_base( pageID, offset)`
- `get_indirection_base(pageld, offset)`
- `get_record_element(pageld, offset, col):`

## Setters:

- `set_indirection_base(pageld, offset, new_indirection)`
- `set_schema_encoding_base(pageld, offset, new_schema_encoding)`

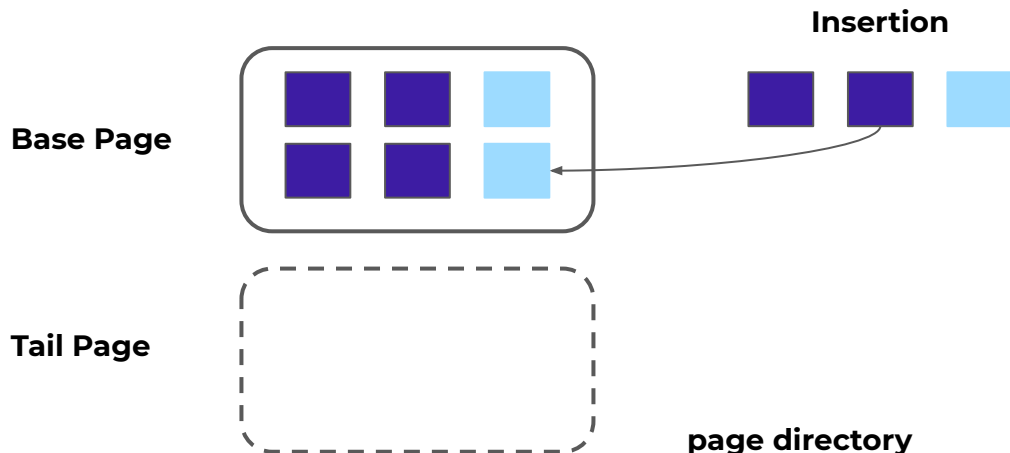
## key\_Exists:

- Check if record exists or has been deleted



# Query Interface: INSERT

INSERT (\*columns)



**page directory**

append ( key = RID, value = (pageID, offset))

**index directory**

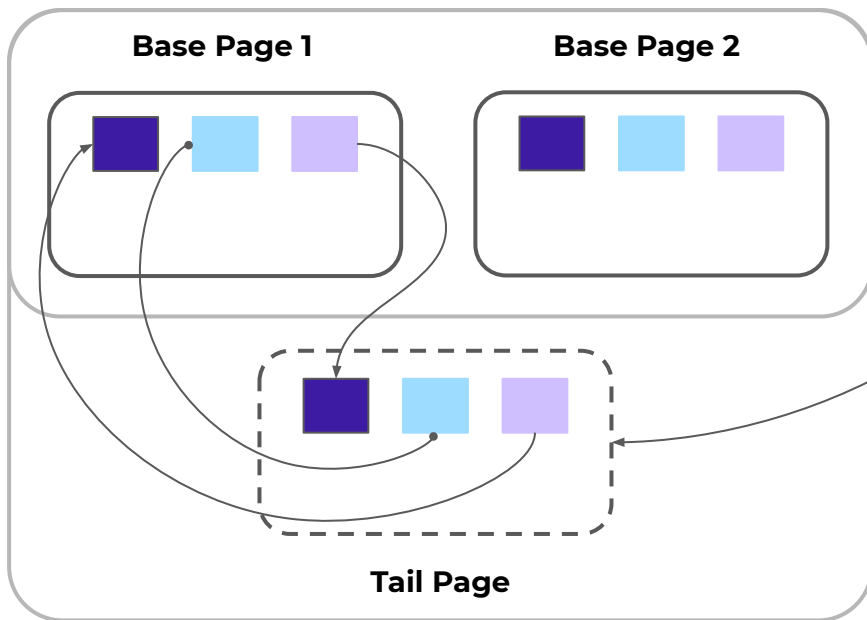
append ( key = KEY, value = RID)



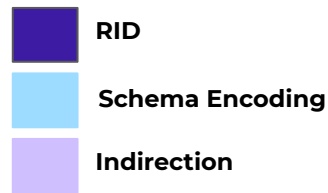
# Query Interface: UPDATE

UPDATE (KEY, \*columns)

Page Range 1



Update on  
Base Page 1



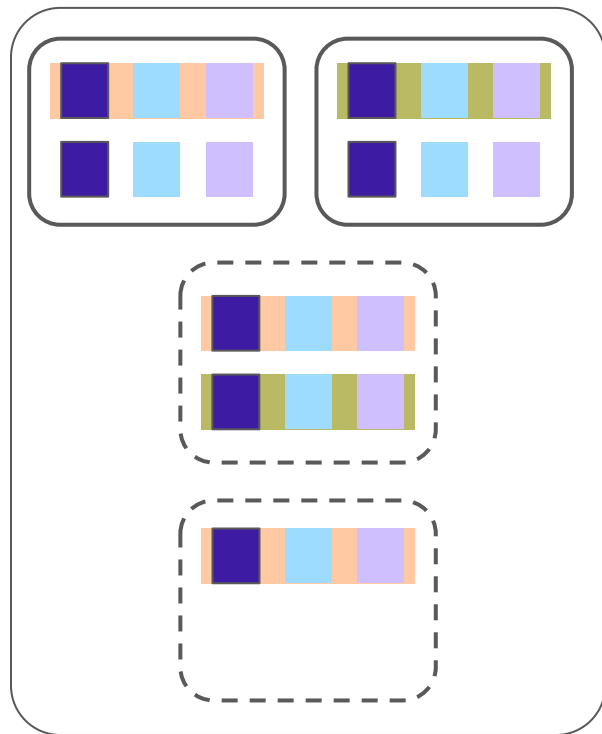
page\_range\_size = 2

# Query Interface: SELECT



SELECT(KEY, column, query\_columns)

RID	Schema Encoding	Indirection	Value1	Value2	Value3
b1	110	t2	9	5	18
t1	010	b1	maxInt	3	maxInt
t2	110	t1	2	3	maxInt





# Query Interface: DELETE, SUM



DELETE (KEY)

<b>Index Directory</b>	<b>RID</b>
<b>Schema_Encoding</b>	<b>'0' * num_columns (for base page)</b>
<b>UPDATE</b>	<b>update(key, *([None] * num_columns))</b>

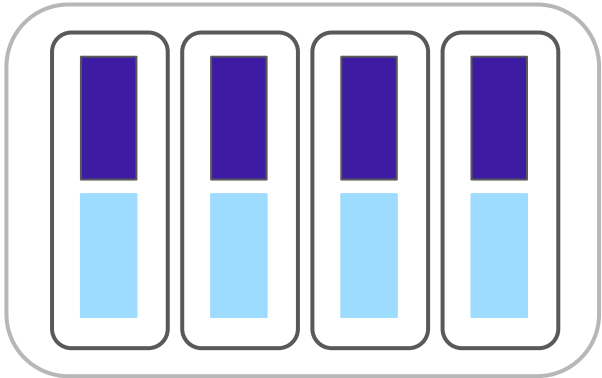
SUM(start\_range, end\_range, aggregate\_column\_index)

<b>key_list</b>	<b>[start_range, ... , end_range]</b>
<b>query_column</b>	<b>query_column[aggregate_column_index] = 1</b>
<b>SUM (SELECT)</b>	<b>sum += select(key, 0, query_column)[0].columns[aggregate_column_index]</b>

# Things to Improve Upon for M2

- Expand query capabilities
  - More functionalities
- Non-int values conversion
- Indexing columns
  - Page Range
  - locate() and locate\_range()
  - B-Tree
- Clean up the CODE

# Alternate Implementation: Base Page and Tail Page in One Page Object



Page Range 1

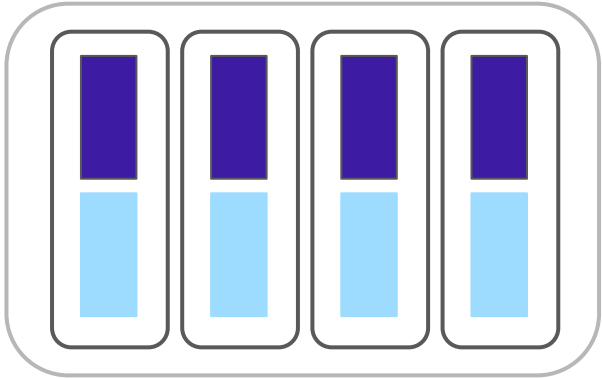
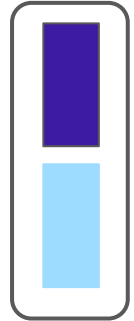
Base Page



Tail Page

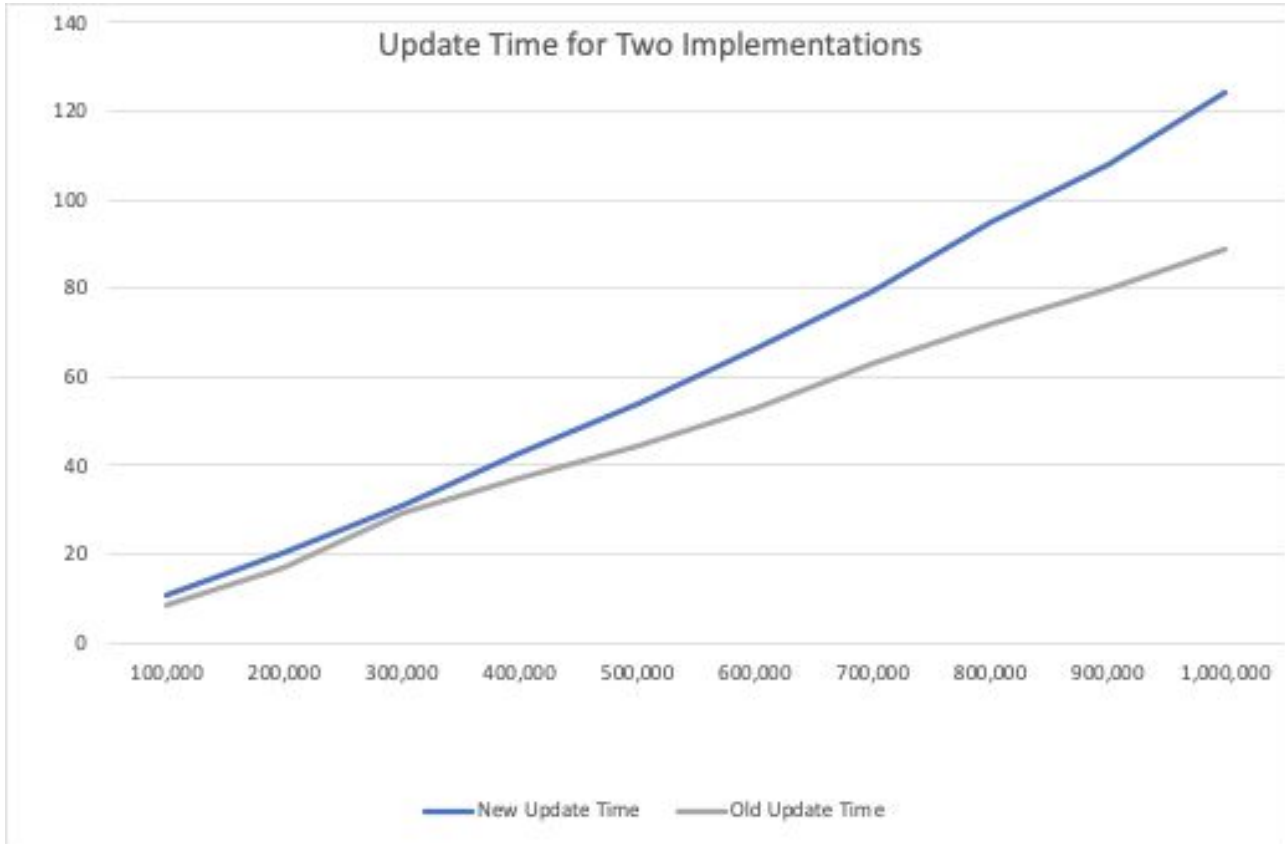


Page Object



Page Range 2

# Alternate Implementation Speed



**Thank You!**