HotStuff: BFT Consensus in the Lens of Blockchain

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Diagrams pulled from:

<u>https://expolab.org/ecs265-fall-2021/slides/4_HotStuff.pdf</u> by Xianda Hou, Oliver Shen, Ashwin Sekhari, Sheshavishnuprasad D, Mythreya K <u>https://expolab.org/ecs265-fall-2022/slides/HotStuff-presentation.pdf</u> by Tong Zhu, Hongxiang Zhang, Siyuan Liu, Yifeng Shi,

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PBFT Overview

HotStuff is based off of a Byzantine Fault Tolerance protocol which allows the distributed system to achieve consensus via a Leader Node and can handle malicious nodes, for both the primary and replicas.

Protocol supports n>=3f+1, where n is total nodes and f is the number of faulty nodes for the system to still work



Problem with PBFT

- BFT was designed around node counts between 4 and 7, however modern blockchain systems require this node count to scale to the thousands, which PBFT is unequipped to handle
- This is in large part due to communication costs in BFT, as BFT has O(n^2) communication complexity and if a view change is necessary, this raises to O(n^3) communication complexity

Protocol	Correct leader	Authenticator complexity Leader failure (view-change)	f leader failures	Responsiveness
DLS [25]	$O(n^4)$	$O(n^4)$	$O(n^4)$	
PBFT [20]	$O(n^2)$	$O(n^3)$	$O(fn^3)$	\checkmark
SBFT [30]	O(n)	$O(n^2)$	$O(fn^2)$	\checkmark
Tendermint [15] / Casper [17]	$O(n^2)$	$O(n^2)$	$O(fn^2)$	
Tendermint [*] / Casper [*]	O(n)	O(n)	O(fn)	
HotStuff	O(n)	O(n)	O(fn)	\checkmark

HotStuff's Proposed Solution

- HotStuff makes use of a one-to-all message structure to address this scaling problem, where the leader alone communicates to n replicas rather than every replica communicating with each other
- This structure places a large amount of power in the leader's hands, so HotStuff utilizes frequent view changes to avoid excess power being given to a single node
- Linear view changes to support frequent view changes
- Threshold Signature to allow for efficient combining of votes from non-primary replicas and for replicas to verify the leader actually received n-f votes



PBFT:

Hotstuff:



Linear View Change

- Linear View Change means that the cost for a new primary to raise a proposal and conduct consensus is equal to the cost for the original primary to do so
- A NextView Interrupt causes the replicas to send their status to the new primary and can happen during



Threshold Signature

With a (k,n)-threshold signature scheme:

- Each replica has its own unique and distinct private key
- Each replica also possesses a common public key
- When a replica sends a message, it signs it using its private key as an ID
- At least k partial signatures can be combined into a final/complete signature
- If a replica receives a complete signature, it can verify it using the public key

Benefits:

- Reduces the number of signatures in consensus
- Reduces the size of the messages being sent
- Results in far less communication per node
- Less verification complexity



HotStuff Protocol Advancements and Key Features

- Linear View Change
- Optimistic Responsiveness
- Expenses incurred by a new leader to drive the protocol towards consensus are comparable to those of the present leader.
- Solves a liveness problem, the hidden lock problem, with the view change protocol by adding a lock-precursor phase.

Mitigating Liveness Issues with Precursor-Lock Round:

Problem:

The hidden lock problem arises when a leader fails to wait for the expiration time (Δ) of a round.

Solution:

- HotStuff introduces a precursor-lock round before the actual lock round to mitigate the hidden lock problem and prevent liveness violations caused by impatient leaders.
- In this precursor-lock round, the leader ensures hearing from 2F+1 participants, enabling it to ascertain the highest lock value proposed (though not necessarily accepted). This mechanism eliminates the necessity of waiting for the maximum Δ expiration time, enhancing system efficiency and responsiveness.

Phases of HotStuff

PREPARE PHASE 1)





Phases of HotStuff

2) Pre-commit Phase:





Phases of HotStuff



In HotStuff, if a replica waits too long for a message, it automatically moves to the next view to maintain progress.

NextView Interrupt

- Waiting for Messages
- Utility Function: nextView(ViewNumber)

HotStuff



Safety

- 1. No 2 conflicting QC will have the same view number. Voting only happens once.
- 2. No 2 conflicting nodes can be committed by correct replicas at any time.
- Atleast 2f + 1 replicas will be non faulty
- Quorum Certificate
- View Changes
- SafeNode Predicate

Liveness

PBFT

- Prevents system from getting stalled
- Avoiding deadlocks

p

ľ1

ľ2

r3

• Optimistically Responsive

Client



What is Chained HotStuff and why do we need it?

Decrease type of messages Only one type of QC ever created

Simplify the protocol Easier to code event-driven style programs

Conceptually separate safety and liveness

View existing BFT protocols in a common framework

Idea : Pipeline the phases



Idea: Pipeline the phases



But L(v2) needs to start its own view...



continues



Same for L(v3), L(v4)



Code is identical for all views



Code for different phases

Prepare Phase:

 Leader sends QC
On receiving QC, check for conflicts and vote Pre-commit Phase:

1. Leader sends QC

2. On receiving QC, lock the node

Commit Phase:

 Leader sends QC
On receiving QC, check for conflicts and vote viewNumber

parent

grand-parent

great-grand-parent

Quorum Certificate for message

Node digests for verification



View = height

One-chain, Two-chain and Three-chains









DLS - Commit after one chain



(a) One-Chain (DLS, 1988)

- Voting node = locked node
- Only leader commits
- Complicated unlock mechanism

PBFT - Commit after two chains



(b) Two-Chain (PBFT, 1999)

- Lock after one phase
- Commit longest chain
- Convincing longest chain requires quadratic messages per view.

Tendermint - Commit after two chains, but wait...



(c) Two-Chain w/ delay (Tendermint, 2016)

- Wait for max delay (after GST) each phase
- This guarantees liveness (no need for large proofs of one-chains)

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