

The Honey Badger of BFT Protocols

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Timing assumptions considered harmful

- Weak Synchrony and Asynchronous BFT Protocol
- Why not weak synchrony
 - The liveness will fail when expected time assumptions are violated
 - Less throughput when network is unpredictable
- HoneyBadgerBFT guarantees liveness without making time assumptions

Asynchronous networks are the “harsh climates” of distributed computing

Full Synchrony: all messages are delivered within Δ time

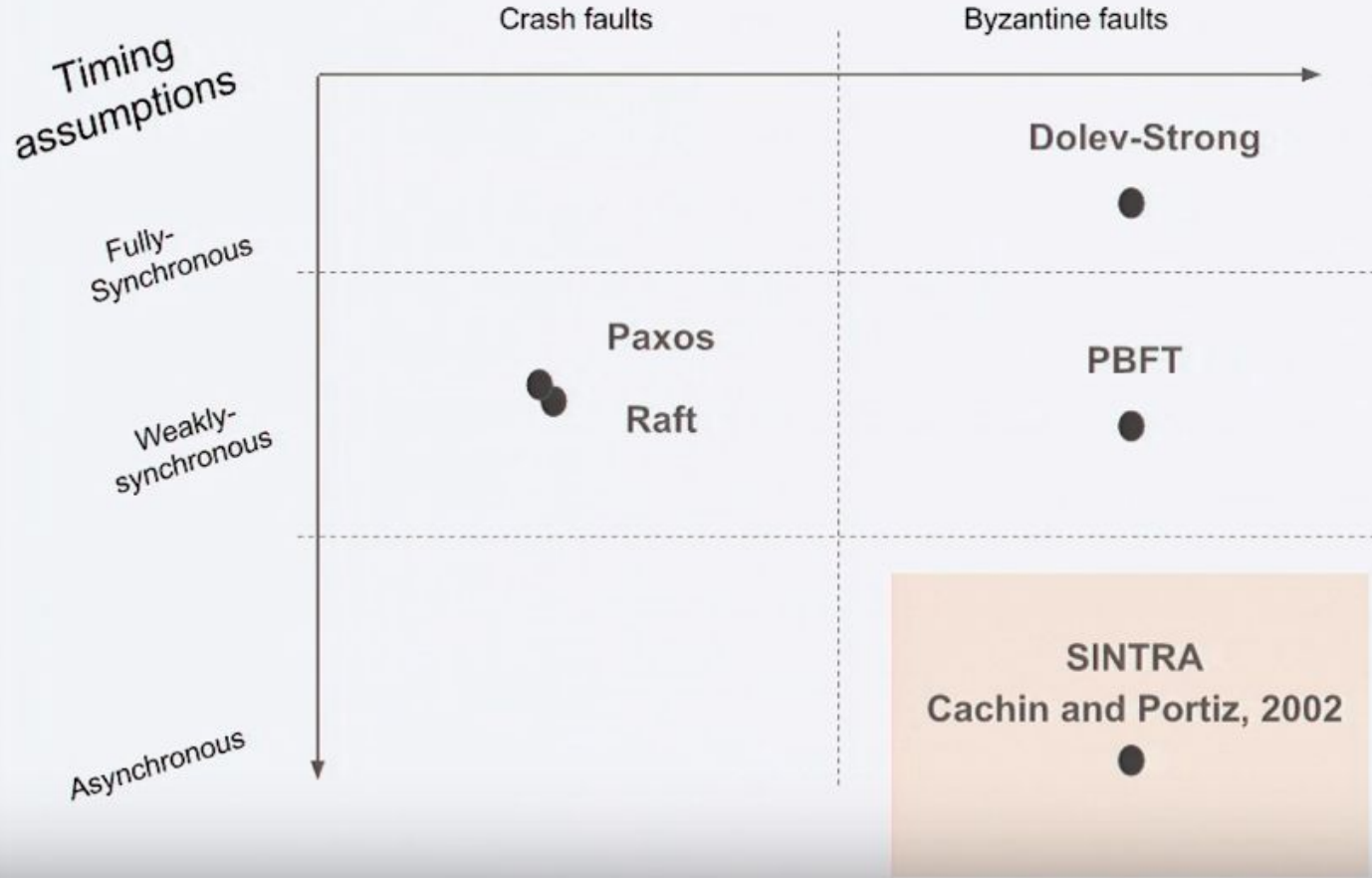
Eventual Synchrony: after unknown time GST , all messages delivered within Δ

Partial Synchrony: Δ is unknown to the protocol

Weak Synchrony: Δt is time varying, but grows polynomially in t

Asynchronous: all messages are eventually delivered

Fault models



The Approach

Adapt synchronous BFT for efficiency in the batch setting

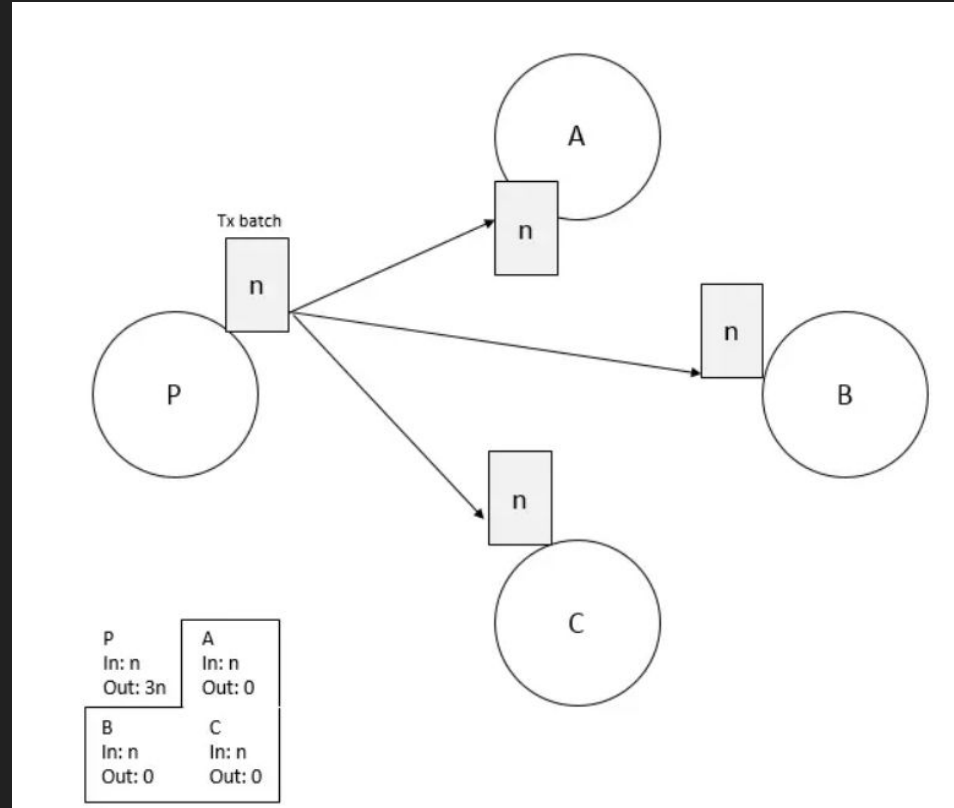
1. Improve by $O(N)$ by “refactoring” with known (but overlooked) primitives
2. Improve by another $O(N)$ by using random selection and threshold encryption

Refactor of the transactions to mitigate node bandwidth bottleneck

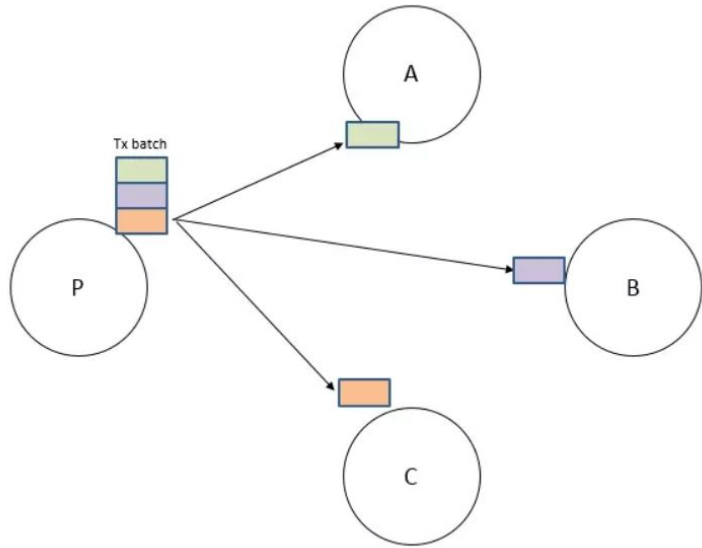
PBFT broadcast standard way

P: leader

Leader bandwidth $O(n*3)$

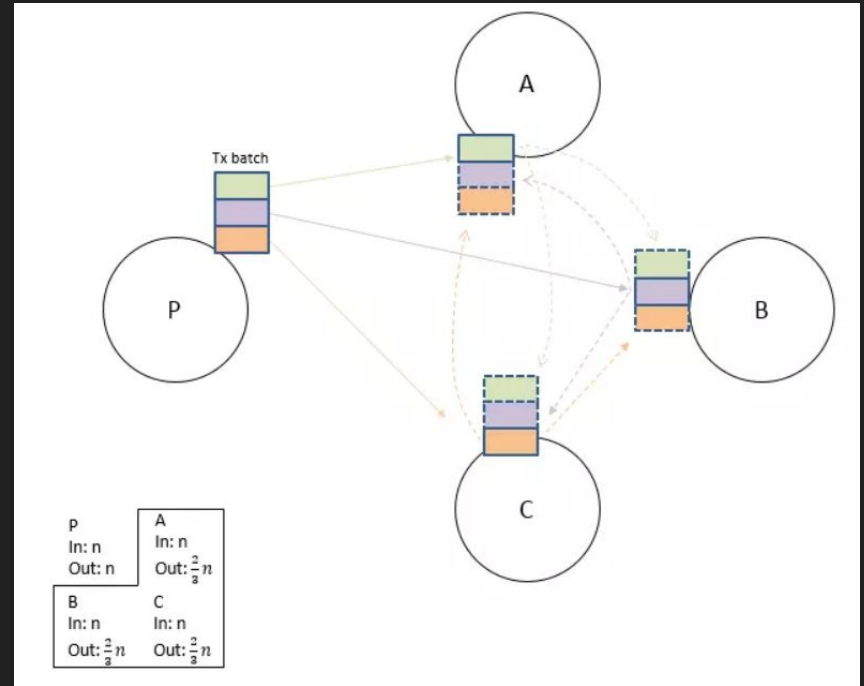


Refactor of the transactions to mitigate node bandwidth bottleneck

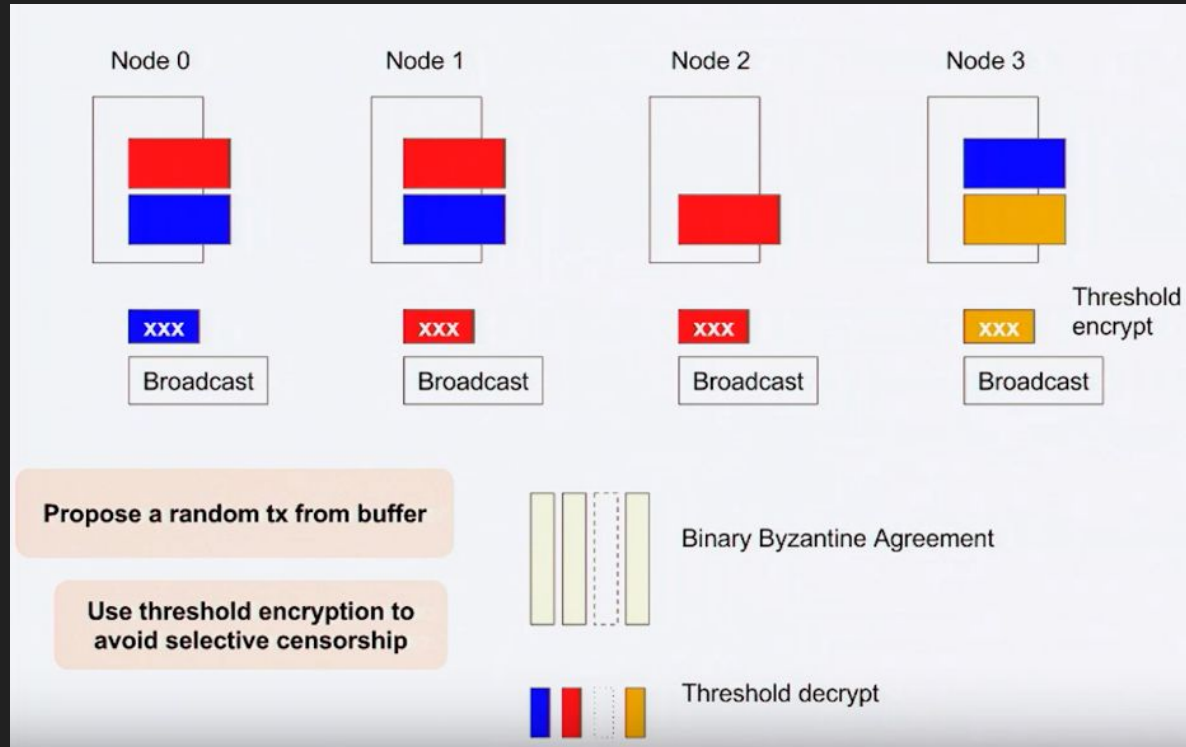


Improved broadcast way

Leader bandwidth $O(n)$



Avoid sending redundant transactions-random selection and threshold encryption



Results: Optimal resilience and efficiency

Choose a large enough batch size, of $B = \Omega(\lambda N^2 \log N)$.

Total Bandwidth per transaction (for each node) is $O(1)$.

Expected # of rounds is $O(\log N)$.

Implementation

Python protocol implementation, using gevent

Threshold cryptography: Charm/PBC library

Signature: Boldyreva '03

Encryptions: Baek and Zhang '03

Experiments on local cluster, worldwide EC2, & over Tor

Thank you!