L-Store: Towards a Unified OLTP and OLAP over a Secure Platform

Mohammad Sadoghi

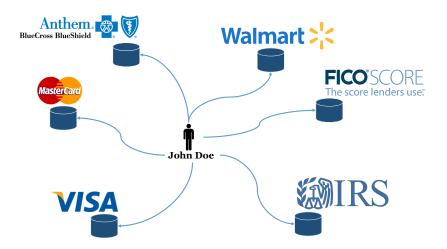
Exploratory Systems Lab University of California, Davis

University of Waterloo October 15, 2018





Data Management at Macroscale

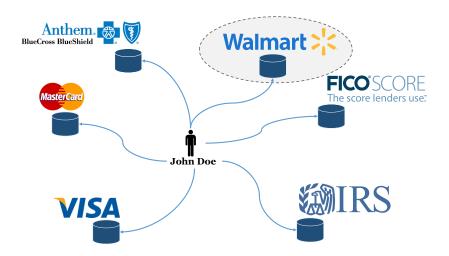


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 L-Store
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 Vision:
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Data Management at Macroscale



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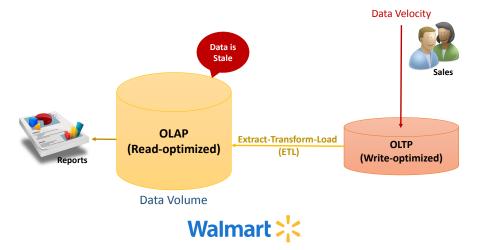
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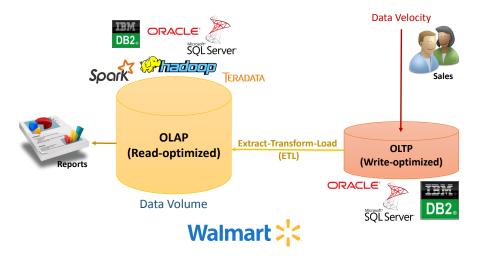




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One Size Does not Fit All As of 2012

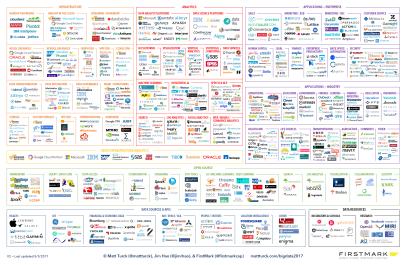
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Big Data Landscape

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One Size Does not Fit All As of 2017

BIG DATA LANDSCAPE 2017



Mohammad Sadoghi (UC Davis)

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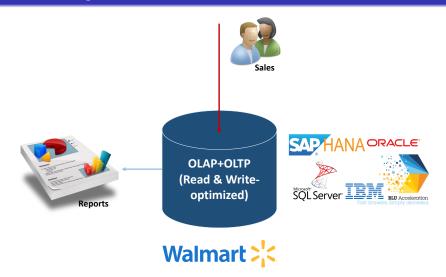
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Data Management at Microscale: Volume & Velocity



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Data Management at Microscale: Volume & Velocity

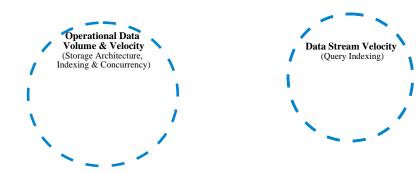


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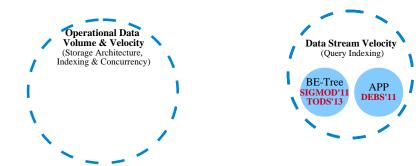
- 2 Data Management at Microscale
- 3 Data Velocity: Index Maintenance
- 4 Data Volume: MVCC Concurrency
- 5 Data Volume: Coordination-free Concurrency
- 6 Combining Volume & Velocity: Lineage-based Storage Architecture
- 7 Data at Macroscale: Decentralized & Democratic Data Platform
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- 9 References





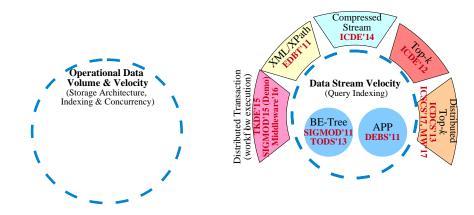
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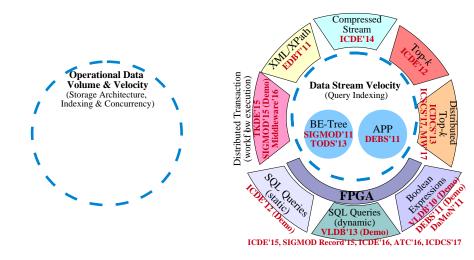




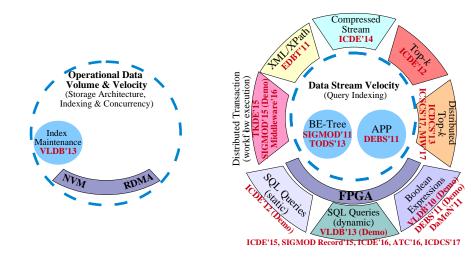
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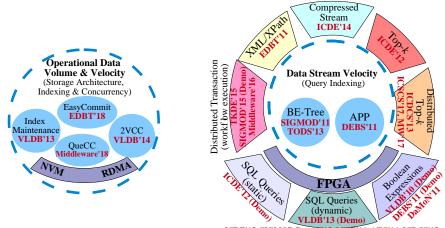






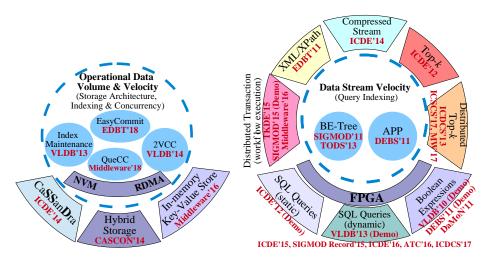




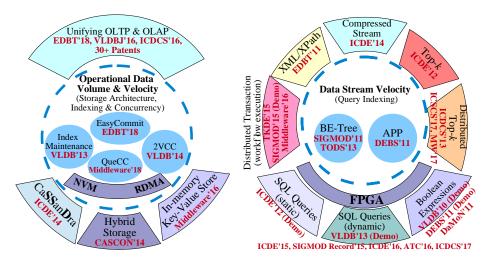


ICDE'15, SIGMOD Record'15, ICDE'16, ATC'16, ICDCS'17





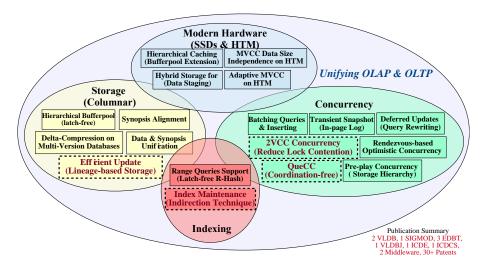




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Deep Dive: Unifying OLTP & OLTP



- 1 Data Management at Microscale
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Extending Storage Hierarchy with Indirection Layer

QueCC

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Evaluation

Vision: L-Store

Conclusions

References



Motivations

Big Picture

Indirection

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Reducing Index maintenance: Velocity Dimension

Observed Trends

In the absence of in-place updates in operational multi-version databases, the cost of index maintenance becomes a major obstacle to cope with data velocity.

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Reducing Index maintenance: Velocity Dimension

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Extending storage hierarchy (using fast non-volatile memory) with *an extra level of indirection* in order to

Reducing Index maintenance: Velocity Dimension

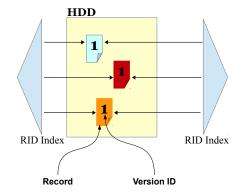
Observed Trends

In the absence of in-place updates in operational multi-version databases, the cost of index maintenance becomes a major obstacle to cope with data velocity.

Extending storage hierarchy (using fast non-volatile memory) with an extra level of indirection in order to Decouple Logical and Physical Locations of Records to Reduce Index Maintenance

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Traditional Multi-version Indexing: Updating Records



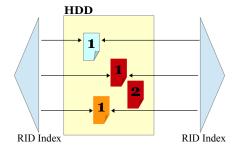
Updating random leaf pages

Mohammad Sa	doghi (UC Davis)
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Traditional Multi-version Indexing: Updating Records

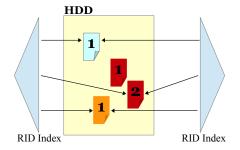


Updating random leaf pages

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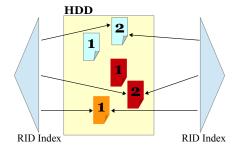
Traditional Multi-version Indexing: Updating Records



Updating random leaf pages

Mohammad S	Sadoghi ((UC D	avis)
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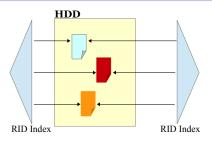
Traditional Multi-version Indexing: Updating Records



Updating random leaf pages

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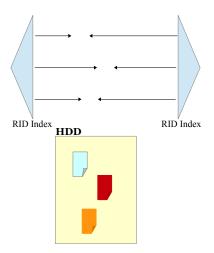
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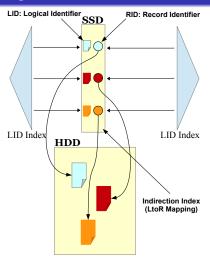
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Indirection Indexing: Updating Records



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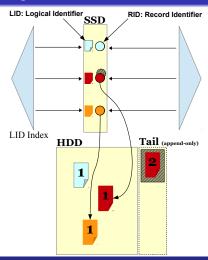
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Indirection Indexing: Updating Records

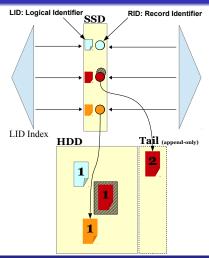


Eliminating random leaf-page updates

Mohammad Sadoghi (UC Davis)

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Indirection Indexing: Updating Records

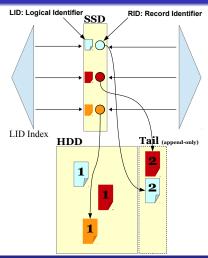


Eliminating random leaf-page updates

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Indirection Indexing: Updating Records



Eliminating random leaf-page updates

Mohammad Sadoghi (UC Davis)

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Analytical & Experimental Evaluations

Image: A matrix

Indirection Time Complexity Analysis

	Legend
K	Number of indexes
LB	LIDBlock size
М	Number of matching records

Method	Туре	Imm. SSD	Def. SSD	Imm. HDD	Def. HDD
Base	Deletion	0	0	2 + K	$\leq 1 + K$
	Single-attr. update	0	0	3 + K	$\leq 2 + K$
	Insertion	0	0	1 + K	$\leq 1 + K$
	Search Uniq.	0	0	2	0
	Search Mult.	0	0	1 + M	0
Indirection	Deletion	2	0	2	≤ 3
	Single-attr. update	2	0	4	≤ 3
	Insertion	2 + 2K	2K/LB	1	$\leq 1 + 2K/LB$
	Search Uniq.	2	0	2	0
	Search Mult.	1 + M	0	1 + M	0

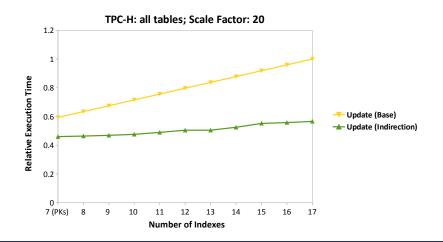
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Experimental Setting

- Hardware:
 - (2 × 8-core) Intel(R) Xeon(R) CPU E7-4820 @ 2.00GHz, 32GB, 2 × HDD, SSD Fusion-io
- Software:
 - Database: IBM DB2 9.7
 - Prototyped in a commercial proprietary database
 - Prototyped in Apache Spark by UC Berkeley
 - LIBGist v.1.0: Generalized Search Tree C++ Library by UC Berkeley (5K LOC) (Predecessor of Generalized Search Tree (GiST) access method for PostgreSQL)
 - LIBGist^{mv} Prototype: Multi-version Generalized Search Tree C++ Library over LIBGist supporting Indirection/LIDBlock/DeltaBlock (3K LOC)
- Data:
 - TPC-H benchmark
 - Microsoft Hekaton micro benchmark



Indirection: Effect of Indexes in Operational Data Stores

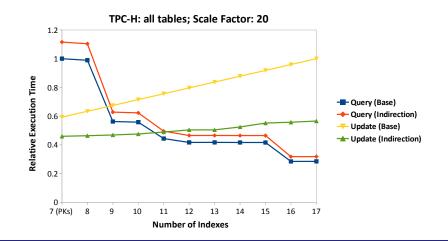


Substantially improving the update time ...

Mohammad Sadoghi (UC Davis)



Indirection: Effect of Indexes in Operational Data Stores



... Consequently affording more indexes and significantly reducing the query time

- 1 Data Management at Microscale
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Introducing Multi-version Concurrency Control

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2VCC



QueCC L-Store

Vision: L-Store

Evaluation

Conclusions

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Motivations

Big Picture Indirection

Generalized Concurrency Control: Volume Dimension

Observed Trends

In operational multi-version databases, there is a tremendous opportunity to avoid clashes between readers (scanning a large volume of data) and writers (frequent updates).

Generalized Concurrency Control: Volume Dimension

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Introducing a (latch-free) *two-version concurrency control (2VCC)* by extending indirection mapping (i.e., central coordination mechanism) and exploiting existing two-phase locking (2PL) in order to

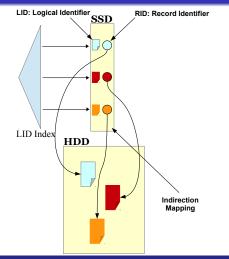
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Introducing a (latch-free) *two-version concurrency control (2VCC)* by extending indirection mapping (i.e., central coordination mechanism) and exploiting existing two-phase locking (2PL) in order to Decouple Readers/Writers to Reduce Contention (Pessimistic and Optimistic Concurrency Control Coexistence)

2V-Indirection Indexing: Updating Records

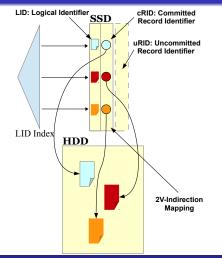


Recap: Indirection technique for reducing index maintenance

Mohammad	Sadoghi ((UC Davis)	
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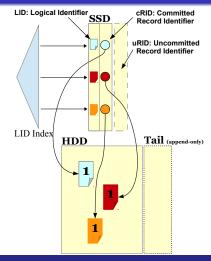
2V-Indirection Indexing: Updating Records



Extending the indirection to committed/uncommitted records

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2V-Indirection Indexing: Updating Records

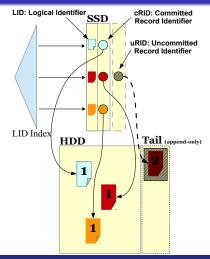


Extending the indirection to committed/uncommitted records

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Image: A matrix

2V-Indirection Indexing: Updating Records

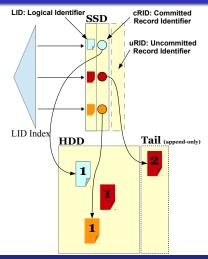


Decoupling readers/writers to eliminate contention

Mohammad Sadoghi (UC Davis)

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2V-Indirection Indexing: Updating Records

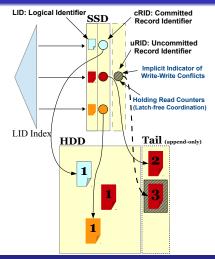


Decoupling readers/writers to eliminate contention

Mohammad	Sadoghi ((UC Davis)	
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2V-Indirection Indexing: Updating Records



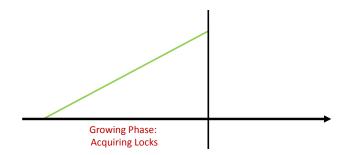
Decoupling readers/writers to eliminate contention

Mohammad Sadoghi (UC Davis)

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Two-phase locking (2PL) consisting of growing and shrinking phases

Overview of Two-version Concurrency Control Protocol

QueCC L-Store

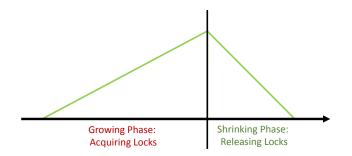
Vision: L-Store

Conclusions

References

Evaluation

2VCC



Two-phase locking (2PL) consisting of growing and shrinking phases

Motivations Big Picture Indirection

Overview of Two-version Concurrency Control Protocol

QueCC

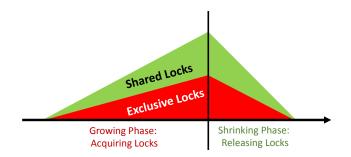
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Two-phase locking (2PL) consisting of growing and shrinking phases

Motivations Big Picture Indirection

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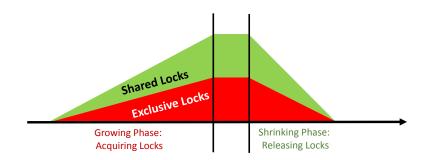
Vision: L-Store

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Evaluation

2VCC



Extending 2PL with certify phase

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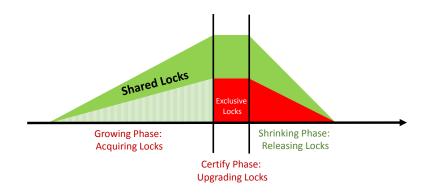
Motivations Big Picture Indirection

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Overview of Two-version Concurrency Control Protocol



Exclusive locks held for shorter period (inherently optimistic)

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QueCC L-Store 00000000 0000000 000000 000000 00000 <u>Overview of Two-version Concurrency Control Protocol</u>

Vision: L-Store

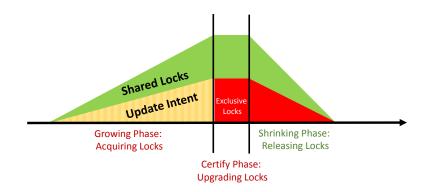
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2VCC



Exclusive locks held for shorter period (inherently optimistic)

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Motivations

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Overview of Two-version Concurrency Control Protocol

Vision: L-Store

Conclusions

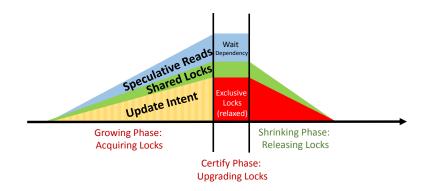
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Indirection

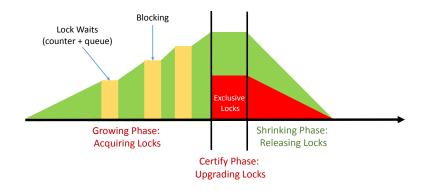


Relaxed exclusive locks to allow speculative reads (increased optimism)

Motivations



Overview of Two-version Concurrency Control Protocol



Trade-offs between blocking (i.e., locks) vs. non-blocking (i.e., read counters)

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Experimental Analysis

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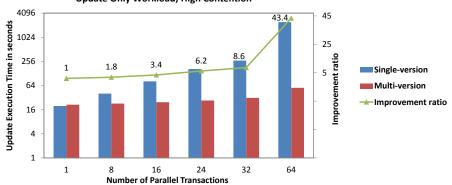
2VCC: Effect of Parallel Update Transactions

2VCC

QueCC

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Indirection



Vision: L-Store

Evaluation

Update Only Workload; High Contention

Substantial gain by reducing the read/write contention & using non-blocking operations

Motivations

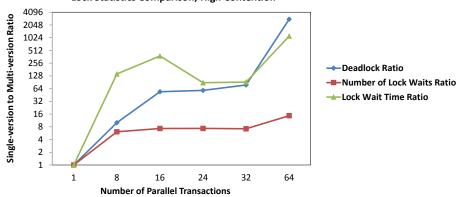
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2VCC: Effect of Parallel Update Transactions



Lock Statistics Comparison; High Contention

Substantial gain by reducing the read/write contention & using non-blocking operations

- 1 Data Management at Microscale
- 2 Data Management at Microscale
- 3 Data Velocity: Index Maintenance
- 4 Data Volume: MVCC Concurrency
- 5 Data Volume: Coordination-free Concurrency
- 6 Combining Volume & Velocity: Lineage-based Storage Architecture
- 7 Data at Macroscale: Decentralized & Democratic Data Platform
- 8 Conclusions
- 9 References

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Introducing Coordination-free Concurrency Control

QueCC

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Vision: L-Store

Evaluation

Conclusions

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Motivations

Big Picture

Indirection

Observed Trends

In operational databases, the use of pre-compiled stored procedures is predominant. There is a tremendous opportunity to exploit transaction prior knowledge to eliminate the need for coordination.

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Observed Trends

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Is it possible to have concurrent execution over shared data (not limited to partitionable workloads) without having any concurrency controls?

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In operational databases, the use of pre-compiled stored procedures is predominant. There is a tremendous opportunity to exploit transaction prior knowledge to eliminate the need for coordination.

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Introducing a *queue-oriented*, *control-free concurrency* (*QueCC*) based on two parallel & independent phases of priority-driven planning & execution.

Observed Trends

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Introducing a *queue-oriented*, *control-free concurrency (QueCC)* based on two parallel & independent phases of priority-driven planning & execution. Execution and Synchronization Decoupling

Queue-oriented, Control-free Concurrency (QueCC)

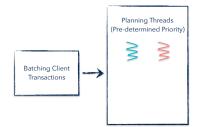
Batching Client Transactions

Execution & Synchronization Decoupling: Deterministic priority-based planning followed by queue-oriented execution

Mohammad Sadoghi (UC Davis)

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Motivations Big Picture Indirection 2VCC QueCC L-Store Evaluation Vision: L-Store Conclusions References coo coocococo cocococo cococo cococo

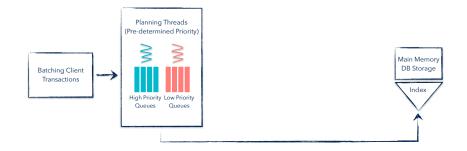


Execution & Synchronization Decoupling: Deterministic priority-based planning followed by queue-oriented execution

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Execution & Synchronization Decoupling: Deterministic priority-based planning followed by queue-oriented execution

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Queue-oriented, Control-free Concurrency (QueCC)

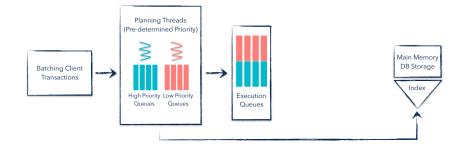
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Vision: L-Store

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Execution & Synchronization Decoupling: Deterministic priority-based planning followed by queue-oriented execution

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Motivations

Big Picture Indirection

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Queue-oriented, Control-free Concurrency (QueCC)

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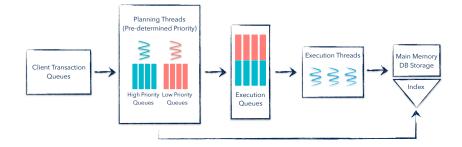
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Conclusions

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Execution & Synchronization Decoupling: Deterministic priority-based planning followed by queue-oriented execution

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Motivations

Big Picture Indirection

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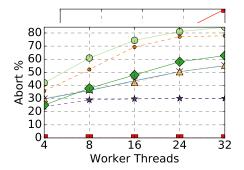
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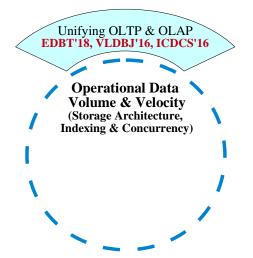
QueCC: Effect of Parallel Update Transactions





Avoiding thread coordination & eliminating all execution-induced aborts 👔 🕤

Unifying OLTP and OLAP



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Unifying OLTP and OLAP: Velocity & Volume Dimensions

Observed Trends

In operational databases, there is a pressing need to close the gap between the write-optimized layout for OLTP (i.e., row-wise) and the read-optimized layout for OLAP (i.e., column-wise).

Image: Image:

Unifying OLTP and OLAP: Velocity & Volume Dimensions

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In operational databases, there is a pressing need to close the gap between the write-optimized layout for OLTP (i.e., row-wise) and the read-optimized layout for OLAP (i.e., column-wise).

Introducing a *lineage-based storage architecture*, a contention-free update mechanism over a native columnar storage in order to

Unifying OLTP and OLAP: Velocity & Volume Dimensions

Observed Trends

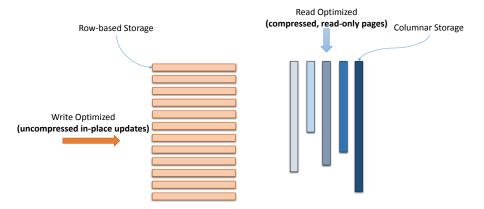
In operational databases, there is a pressing need to close the gap between the write-optimized layout for OLTP (i.e., row-wise) and the read-optimized layout for OLAP (i.e., column-wise).

Introducing a *lineage-based storage architecture*, a contention-free update mechanism over a native columnar storage in order to

lazily and independently stage stable data from a write-optimized layout (i.e., OLTP) into a read-optimized layout (i.e., OLAP)

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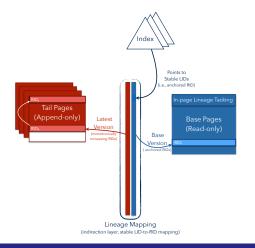
Storage Layout Conflict



Write-optimized (i.e., uncompressed & row-based) vs. read-optimized (i.e., compressed & column-based) layouts

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Lineage-based Storage Architecture (LSA): Intuition



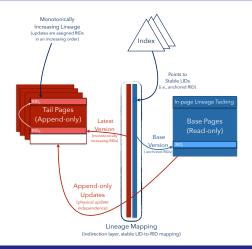
Physical Update Independence: De-coupling data & its updates (reconstruction via in-page lineage tracking and lineage mapping)

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Lineage-based Storage Architecture (LSA): Intuition



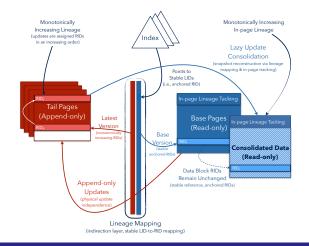
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Lineage-based Storage Architecture (LSA): Intuition



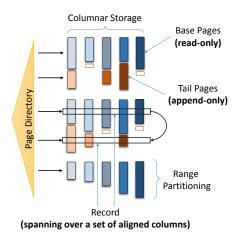
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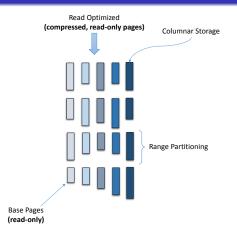
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Lineage-based Storage Architecture (LSA): Overview



Overview of the lineage-based storage architecture (base pages and tail pages are handled identically at the storage layer)

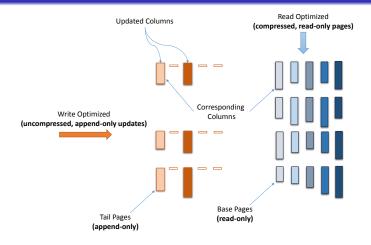
L-Store: Detailed Design



Records are range-partitioned and compressed into a set of ready-only **base pages** (accelerating analytical queries)

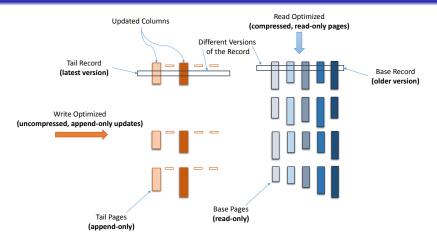
Mohammad	Sadoghi ((UC Davis)	
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L-Store: Detailed Design



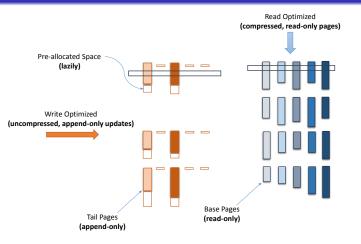
Recent updates for a range of records are clustered in their **tails pages** (transforming costly point updates into an amortized analytical-like query)

L-Store: Detailed Design



Recent updates for a range of records are clustered in their **tails pages** (transforming costly point updates into an amortized analytical-like query)

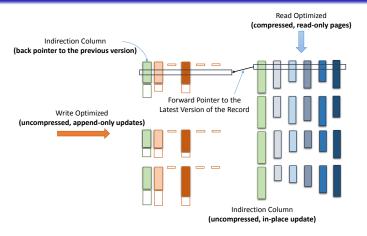
L-Store: Detailed Design



Recent updates are strictly appended, uncompressed in the pre-allocated space (eliminating the read/write contention)

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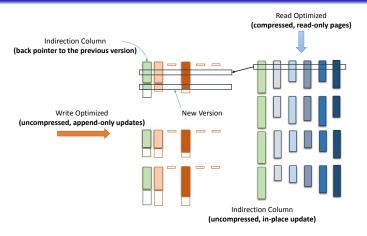
L-Store: Detailed Design



Achieving (at most) 2-hop access to the latest version of any record (avoiding read performance deterioration for point queries)

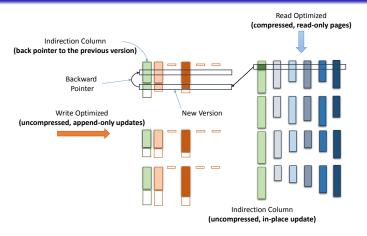
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L-Store: Detailed Design



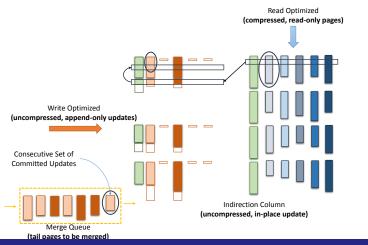
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L-Store: Detailed Design



Achieving (at most) 2-hop access to the latest version of any record (avoiding read performance deterioration for point queries)

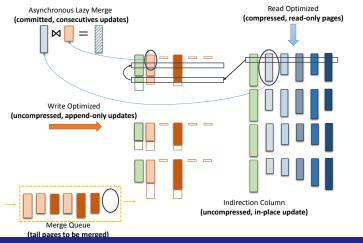
L-Store: Contention-free Merge



Contention-free merging of only stable data: read-only and committed data (no need to block on-going and new transactions)

Mohammad Sadoghi (UC Davis)

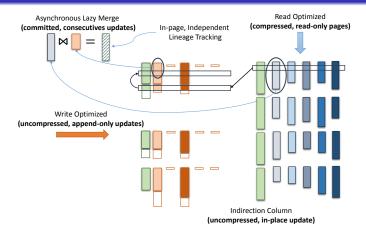
L-Store: Contention-free Merge



Lazy independent merging of **base pages** with their corresponding **tail pages** (resembling a local left outer-join of the base and tail pages)

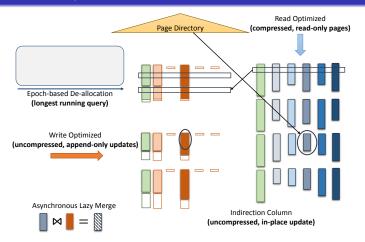
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L-Store: Contention-free Merge



Independently tracking the lineage information within every page (no need to coordinate merges among different columns of the same records)

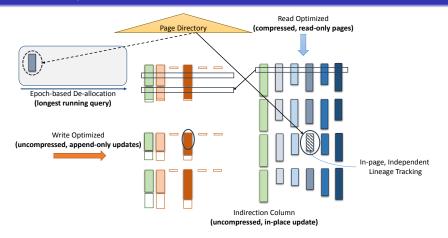
L-Store: Epoch-based Contention-free De-allocation



Contention-free page de-allocation using an epoch-based approach (no need to drain the ongoing transactions)

Mohammad Sadoghi (UC Davis)

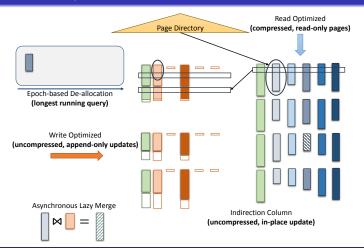
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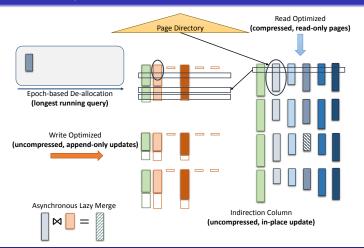
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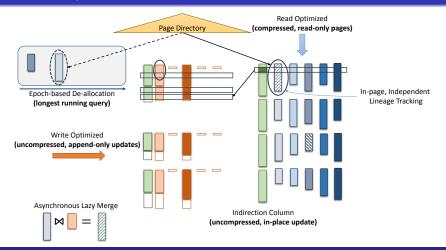
L-Store: Epoch-based Contention-free De-allocation



Contention-free page de-allocation using an epoch-based approach (no need to drain the ongoing transactions)

Mohammad Sadoghi (UC Davis)

L-Store: Epoch-based Contention-free De-allocation



Contention-free page de-allocation using an epoch-based approach (no need to drain the ongoing transactions)

Experimental Analysis

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Experimental Settings

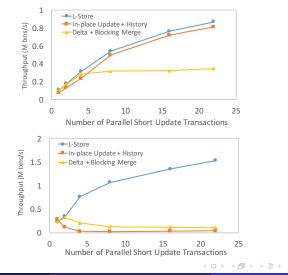
Hardware:

• 2 \times 6-core Intel(R) Xeon(R) CPU E5-2430 @ 2.20GHz, 64GB, 15 MB L3 cache

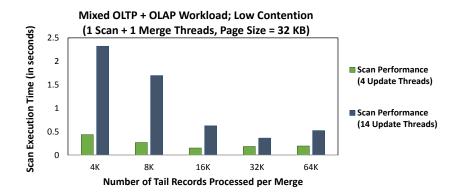
Workload: Extended Microsoft Hekaton Benchmark

- Comparison with In-place Update + History and Delta + Blocking Merge
- Effect of varying contention levels
- Effect of varying the read/write ratio of short update transactions
- Effect of merge frequency on scan
- Effect of varying the number of short update vs. long read-only transactions
- Effect of varying L-Store data layouts (row vs. columnar)
- Effect of varying the percentage of columns read in point queries
- Comparison with log-structured storage architecture (*LevelDB*)

Effect of Varying Contention Levels

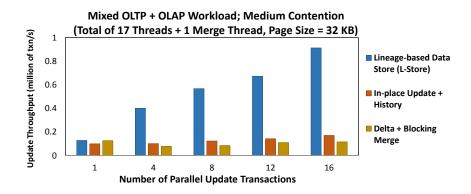


Effect of Merge Frequency on Scan Performance



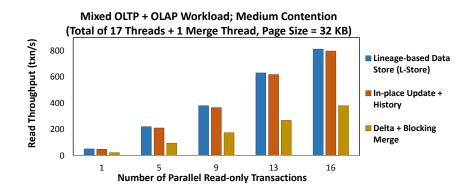
Merge process is essential in maintaining efficient scan performance

Effect of Mixed Workloads: Update Performance



Eliminating latching & locking results in a substantial performance improvement

Effect of Mixed Workloads: Read Performance



Coping with tens of update threads with a single merge thread

- 1 Data Management at Microscale
- 2 Data Management at Microscale
- 3 Data Velocity: Index Maintenance
- 4 Data Volume: MVCC Concurrency
- 5 Data Volume: Coordination-free Concurrency
- 6 Combining Volume & Velocity: Lineage-based Storage Architecture
- 7 Data at Macroscale: Decentralized & Democratic Data Platform
- 8 Conclusions
- 9 References

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Recap: Data Management Challenges at Microscale



OLTP and OLAP data are isolated at microscale

Mohammad Sadoghi (UC Davis)

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Recap: Data Management Challenges at Microscale



First step is to unify OLTP and OLAP

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Platform Scaling: Data Partitioning

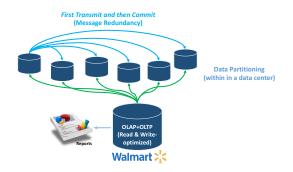


Moving towards distributed environment

Mohammad Sa	doghi (UC Davis)
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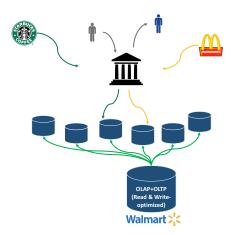
Platform Scaling: Non-blocking Agreement Protocols



Message redundancy vs. latency trade-offs [EasyCommit, EDBT'18]

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Central Control: Data Gate Keeper



Conform to trusting the central authority and governance

Decentralized Control: Removing Data Barrier

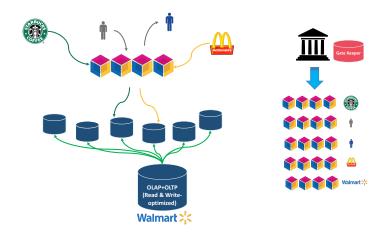


Seek trust in *decentralized* and *democratic* governance [PoE (under submission)]

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Democratic Control: Removing Trust Barrier



Seek trust in *decentralized* and *democratic* governance [PoE (under submission)]

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Global-scale Reliable Platform over Unreliable Hardware

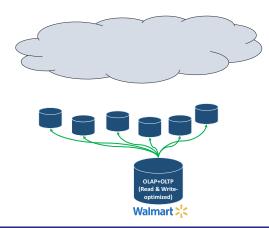


Self-managed infrastructure

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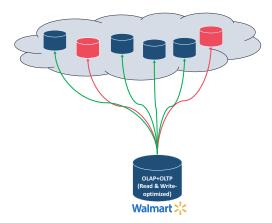
Global-scale Reliable Platform over Unreliable Hardware



Cloud-managed infrastructure (trust the provider)

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Global-scale Reliable Platform over Unreliable Hardware

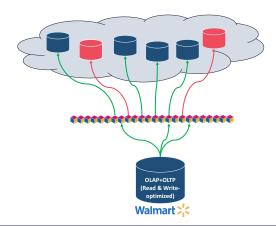


Cloud-managed infrastructure (trust the provider)

Mohammad Sadoghi (UC Davis)
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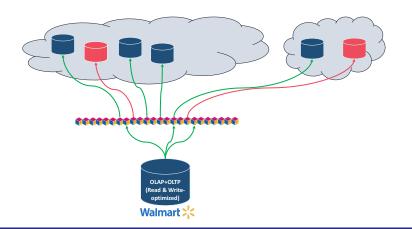
Global-scale Reliable Platform over Unreliable Hardware



Light-weight, fault-tolerant, trusted middleware [Blockplane, (under submission)]

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Global-scale Reliable Platform over Unreliable Hardware

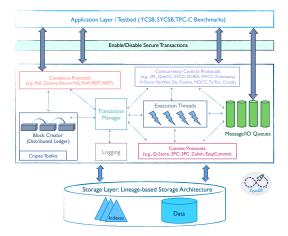


Fault-tolerant protocols vs. consistency models [MultiBFT, GeoBFT (under submission)]

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ExpoDB: Exploratory Data Platform Architecture



A decentralized & democratic platform to unify OLTP and OLAP

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- 1 Data Management at Microscale
- 2 Data Management at Microscale
- 3 Data Velocity: Index Maintenance
- 4 Data Volume: MVCC Concurrency
- 5 Data Volume: Coordination-free Concurrency
- 6 Combining Volume & Velocity: Lineage-based Storage Architecture
- 7 Data at Macroscale: Decentralized & Democratic Data Platform
- 8 Conclusions
- 9 References

Contributions & Outlook

ExpoDB: Decentralized & Democratic Platform

- Decentralized & Democratic Control: PoE, MultiBFT, GeoBFT [under submission]
- Reliability over Unreliable Hardware: Blockplane [under submission]

Operational Data Stores: Velocity & Volume

- Index Maintenance: Indirection Technique [VLDB'13, VLDBJ'16]
- Concurrency Control: 2VCC Technique [VLDB'14, Middleware'16], EasyCommit [EDBT'18], QueCC [Middleware'18]
- Hybrid Storage: Enhancing Key-Value Store [VLDB'12, ICDE'14]
- Real-time OLTP+OLAP: Lineage-based Data Store (L-Store) [EDBT-18, ICDCS'16, 30+ Patents]

Stream Processing: Velocity

- High-dimensional Indexing: BE-Tree [SIGMOD'11, TODS'13], Compressed Stream Processing [ICDE'14]
- (Distributed) Top-k Indexing: BE*-Tree [ICDE'12, ICDCS'13, Middleware'17, ICDCS'17]
- Hardware Acceleration: FPGAs [VLDB'10, ICDE'12, VLDB'13, ICDE'15, SIGMOD Record'15, ICDE'16, USENIX ATC'16, ICDCS'17, ICDE'18]
- Novel Mappings: XML/XPath [EDBT'11], Distributed Workflow [TDKE'15, SIGMOD'15, ICDE'16, Middleware'16]

Questions? Thank you!

Exploratory Systems Lab (ExpoLab) Website: https://msadoghi.github.io/





Related Publications (Patents Omitted)

- 1. Abdelante, A. Folson, O. Hanarmatele, P. Zhang, M. Sarlight Large-scale viscotrant and tentical visibility based mixing of boundings graph to predict drug drug interaction D A. Palean O. Hessenselsh M. Salashi and F. Zhana nnen, G. Franzenzen, in., anzige, zur F. Anzij. Bring deg deg bienation: Wengh similarity kand liek prolition vor ark data. Soundige of the 20th International Conference on World Web Web, WWW 2020, Montred, Canada, April 32 20, 2020, Companies Volume, pages 179–178, 2020 A. Polson, M. Badaghi, G. Hanamashh, and P. Zhang. indiving dog bination though large sole similarly hand lob predation. Expert R. Evolution and M. Enlaght Antiputing database methods by tellisum hardware system on design. or Data Engineering, ICDE 2028, Helsold, Finland, May 20-20, 2028, pages 1428-1411, 2028. 5 A Chandel C. Homomorphic N. Knodes, M. Sashashi and C. Schusters, P & Research M Reducts and M & Decision Insuch colorability hard intensis density with root proming a Propulsion of the 10th ACM intensity or density or Date-Sarah D C. Henneschi, M. Sadarhi and K. J. Mi M. Herendyner, E. Montrashin, M. Reisselerge, and M. Sasinghi. Kanal: A distributed in summer brought stars. D 10 Junior M Redeald and M.A. Juniors M. Joyler, M. Sanlegh, and H. A. Jacoban. DOBOMI: A management infrastructure for distributed data sensitie merkfluor. In Proceedings of the 2021 ACM ISINDIS International Conference on Management and Management a D M. Joyle M. Radarki and H.A. Jacobser. P. Menne, T. Rald, M. Bashaghi, and H. Jamitson. D I Have T Red M Redents and R Jackson D & Mathematica V New M Statistics and M.A. Incoheren D 10 North M. Samuel and M. Inches
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 - M. Najal, M. Sadaghi, and H. Jaminon. The FOP vision: Findlet many around as an around avoid a computing fairly.
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- T. Rall, M. Balleghi, S. Garree-Wilamor, V. Musick-Mulees, H. A. Janubere, and S. Mashookit.
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- M. Sadaghi, M. Carter, R. Bhattacharjee, F. Nagel, and K. A. Rose
- D M. Salashi and H. January
- or Data Engineering, Chicago, NDN 2014, IL, USA, March IV April 6, 2014, march 201-179, 2014
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- D. M. Sadaghi, S. Ebstacherjee, E. Battacharjee, and M. Canin. Lillion: A multime OLTP and OLAP system.
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- M. Sadaghi, H. Singh, and H.-A. Janshon. Transh highly multi-second promition from the second provider bardware.
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- R V. Doors M. Radachi V. Mathematic and H. A. Jacobser.
- B K. Dave M. Sadarbi and H.A. Jacobert K. Zhong M. Sahighi, and H. A. Jarohen. Ek share: A distribution hybrid SCIT and GLAP data proceeding region. Ek share: A distribution hybrid SCIT are not investigated commuting lipsings, SCICI 2018, Raws, Apper, Am 27,30, 2018, pages 309–770, 2018.
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