



## Resilient DB: Building Global-Scale Privacy-Preserving Blockchain Fabric

## Mohammad Sadoghi

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#### **Mohammad Sadoghi**

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Jelle Hellings, PostDoc (Fault-tolerant Complexity Analysis)



Suyash Gupta, PhD (Scalable Consensus Meta-Protocols)



Thamir Qadah, PhD (Distributed & Coordination-free Concurrency)



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Priya Holani, MSc (Scaling Fabric via Sharding)



Shubham Pandey, MSc (Scaling Fabric via RDMA)



Rohan Sogani, MSc (Scaling Fabric via Sharding)





Xinyuan Sun, BSc (Scaling Fabric via RDMA)



FPGA Acceleration: FQP (Flexible Query Processor) [VLDB'10, ICDE'12, VLDB'13, ICDE'15, SIGMOD Record'15, ICDE'16, USENIX ATC'16, ICDCS'17, ICDE'18, TKDE'19]

SQL

**Analytics** 



**SQL** Transactions **SQL** Analytics

High-dimensional Indexing: (e.g., BE-Tree, BE-topK) [SIGMOD'11, ICDE'12, TODS'13, ICDCS'13, ICDE'14, ICDCS'17, Middleware'17]



Concurrency Control Protocols: (e.g., 2VCC, QueCC - Best Paper Award) [VLDB'13, VLDB'14, VLDBJ'16, Middleware'16, TDKE'15, SIGMOD'15, ICDE'16, Middleware'18]

#### QueCC: Queue-Oriented Planning and Execution Architecture





Concurrency Control Protocols: (e.g., 2VCC, QueCC - Best Paper Award) [VLDB'13, VLDB'14, VLDBJ'16, Middleware'16, TDKE'15, SIGMOD'15, ICDE'16, Middleware'18]



[VLDB'12, ICDE'14, ICDCS'16, EDBT'18, 34 filed US patents]

Graphs on SQL: (e.g., GRFusion) [SIGMOD'18, EDBT'18] 7





Consensus Protocols: (e.g., ResilientDB, GeoBFT, PoE, MBFT, Delayed Replication, CSP, Blockplane) [VLDB'20, ICDCS'20, ICDT'20, DISC'19 (2x), SC'19, ICDE'19, arXiv'19 (6x)]

![](_page_9_Picture_0.jpeg)

![](_page_9_Picture_1.jpeg)

![](_page_9_Picture_2.jpeg)

#### **Transaction Processing on Modern Hardware.**

Synthesis Lectures on Data Management, Morgan & Claypool Publishers 2019

#### Fault-Tolerant Distributed Transactions on Blockchain.

Synthesis Lectures on Data Management, Morgan & Claypool Publishers, to appear 2020

![](_page_9_Picture_7.jpeg)

![](_page_9_Picture_8.jpeg)

![](_page_10_Picture_0.jpeg)

![](_page_10_Picture_1.jpeg)

Advancements TV With Ted Danson - CNBC, CityAM, Medium, Yahoo! Finance, Market Insider, CoinDesk, Crypto Media, Davis Enterprise, Times Union, WBOC TV/Radio

### **Books** Transaction Processing on Modern Hardware.

Synthesis Lectures on Data Management, Morgan & Claypool Publishers 2019

#### **Fault-Tolerant Distributed Transactions on Blockchain.**

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![](_page_10_Picture_7.jpeg)

![](_page_10_Picture_8.jpeg)

# **ExpoDB Architecture**

![](_page_11_Figure_1.jpeg)

![](_page_12_Picture_0.jpeg)

# **Quantifiable Resiliency** (Graduate Student Experiments)

Aloha Lake, Desolation Wilderness 15 Miles Long 2,500 Feet Elevation Gain (8,700 Feet at Summit)

![](_page_14_Picture_1.jpeg)

![](_page_14_Picture_2.jpeg)

![](_page_14_Picture_3.jpeg)

#### Tomales Point Trail, Point Reyes National Seashore 9.4 Miles Long 1,579 Feet Elevation Gain

![](_page_15_Picture_1.jpeg)

![](_page_15_Picture_2.jpeg)

# Non-Quantifiable Resiliency

## **Proof-of-Execution: Reaching Consensus Through Fault-Tolerant Speculation [arXiv'19]**

Out-of-Order message processing to reduce replica idleness Speculative Execution with revertible/divergent replicas & eager/irrevertible client commit

![](_page_17_Figure_2.jpeg)

**Fault-tolerant Proof-of-Execution Protocol** 

![](_page_17_Figure_4.jpeg)

PoE scales beyond <u>32 replicas</u>, in presence of failures, outperforms <u>PBFT up to 40%</u>

## **Proof-of-Execution: Reaching Consensus Through Fault-Tolerant Speculation [arXiv'19]**

Out-of-Order message processing to reduce replica idleness Speculative Execution with revertible/divergent replicas & eager/irrevertible client commit

introducing linear message complexity

![](_page_18_Figure_3.jpeg)

**Linearized Proof-of-Execution Protocol** 

## MultiBFT: Scaling Blockchain Databases Through Parallel Resilient Consensus Paradigm [arXiv'19]

A wait-free meta-protocol...

**Designate multiple replicas as Primaries!** 

Run multiple parallel consensuses on each replica independently

![](_page_19_Figure_4.jpeg)

#### **Fault-tolerant MultiBFT Protocol**

20

## MultiBFT: Scaling Blockchain Databases Through Parallel Resilient Consensus Paradigm

![](_page_20_Figure_1.jpeg)

# GeoBFT: Global Scale Resilient Blockchain Fabric [VLDB'20]

A meta-protocol, locally running any BFT in parallel and independently Global ordering provably requires only linear communication Provably sufficient for primary to send a certificate to at most f+1 replicas, malicious primary is detectable and replaceable

![](_page_21_Figure_2.jpeg)

**Fault-tolerant GeoBFT Protocol** 

# GeoBFT: Global Scale Resilient Blockchain Fabric [VLDB'20]

![](_page_22_Figure_1.jpeg)

ResilientDB easily scales across <u>6 countries</u> in <u>4 continents</u> due to GeoBFT protocol.

GeoBFT scales a permissioned blockchain up to <u>60 replicas globally</u>.

# The Fault-Tolerant Cluster-Sending Problem [DISC'19]

formalizing the problem of sending a message from one Byzantine cluster to

another Byzantine cluster in a reliable manner,

establishing lower bounds on the complexity

of this problem under crash failures and Byzantine failures

(linear in the size of clusters)

![](_page_23_Figure_6.jpeg)

# The Fault-Tolerant Cluster-Sending Problem [DISC'19]

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(linear in the size of clusters)

![](_page_24_Figure_6.jpeg)

		Protocol	System	Robustness	Messages	Message size
	ar	RB-bcs	Omit	$\mathbf{n}_{\mathcal{C}_1} > 2\mathbf{f}_{\mathcal{C}_1},  \mathbf{n}_{\mathcal{C}_2} > \mathbf{f}_{\mathcal{C}_2}$	$(\mathbf{f}_{\mathcal{C}_1} + 1) \cdot (\mathbf{f}_{\mathcal{C}_2} + 1)$	$\mathcal{O}(\ v\ )$
_	ine	RB-brs	Byzantine, RS	$\mathbf{n}_{\mathcal{C}_1} > 2\mathbf{f}_{\mathcal{C}_1},  \mathbf{n}_{\mathcal{C}_2} > \mathbf{f}_{\mathcal{C}_2}$	$(2\mathbf{f}_{\mathcal{C}_1}+1)\cdot(\mathbf{f}_{\mathcal{C}_2}+1)$	$\mathcal{O}(\ v\ )$
	h-l	RB-bcs	Byzantine, RS	$\mathbf{n}_{\mathcal{C}_1} > 2\mathbf{f}_{\mathcal{C}_1},  \mathbf{n}_{\mathcal{C}_2} > \mathbf{f}_{\mathcal{C}_2}$	$(\mathbf{f}_{\mathcal{C}_1} + 1) \cdot (\mathbf{f}_{\mathcal{C}_2} + 1)$	$\mathcal{O}(\ v\  + \mathbf{f}_{\mathcal{C}_1})$
	ou	RB-bcs	Byzantine, CS	$\mathbf{n}_{\mathcal{C}_1} > 2\mathbf{f}_{\mathcal{C}_1},  \mathbf{n}_{\mathcal{C}_2} > \mathbf{f}_{\mathcal{C}_2}$	$(\mathbf{f}_{\mathcal{C}_1}+1)\cdot(\mathbf{f}_{\mathcal{C}_2}+1)$	$\mathcal{O}(\ v\ )$
		PBS-bcs	Omit	$\mathbf{n}_{\mathcal{C}_1} > 3\mathbf{f}_{\mathcal{C}_1},  \mathbf{n}_{\mathcal{C}_2} > 3\mathbf{f}_{\mathcal{C}_2}$	$\mathcal{O}(\max(\mathbf{n}_{\mathcal{C}_1},\mathbf{n}_{\mathcal{C}_2})) \text{ (optimal)}$	$\mathcal{O}(\ v\ )$
	ear	PBS-brs	Byzantine, RS	$\mathbf{n}_{\mathcal{C}_1} > 4\mathbf{f}_{\mathcal{C}_1},  \mathbf{n}_{\mathcal{C}_2} > 4\mathbf{f}_{\mathcal{C}_2}$	$\mathcal{O}(\max(\mathbf{n}_{\mathcal{C}_1},\mathbf{n}_{\mathcal{C}_2})) \text{ (optimal)}$	$\mathcal{O}(\ v\ )$
	line	PBS-bcs	Byzantine, RS	$\mathbf{n}_{\mathcal{C}_1} > 3\mathbf{f}_{\mathcal{C}_1},  \mathbf{n}_{\mathcal{C}_2} > 3\mathbf{f}_{\mathcal{C}_2}$	$\mathcal{O}(\max(\mathbf{n}_{\mathcal{C}_1},\mathbf{n}_{\mathcal{C}_2}))$	$\mathcal{O}(\ v\  + \mathbf{f}_{\mathcal{C}_1})$
		PBS-bcs	Byzantine, CS	$\mathbf{n}_{\mathcal{C}_1} > 3\mathbf{f}_{\mathcal{C}_1},  \mathbf{n}_{\mathcal{C}_2} > 3\mathbf{f}_{\mathcal{C}_2}$	$\mathcal{O}(\max(\mathbf{n}_{\mathcal{C}_1},\mathbf{n}_{\mathcal{C}_2}))$ (optimal)	$\mathcal{O}(\ v\ )$

Brief Announcement: The Fault-Tolerant Cluster-sending Problem. DISC 2019

# **Byzantine Cluster-Sending in Expected Constant Communication [arXiv'20]**

formalizing the problem of probabilistically sending a message from one

Byzantine cluster to another Byzantine cluster in a reliable manner,

establishing lower bounds on the complexity

of this problem under crash failures and Byzantine failures

![](_page_25_Figure_5.jpeg)

(expected constant message complexity)

J	Protocol	Robustness	Mes	sage Steps	О.	U.
			(expected)	(worst)		
	PBS-cs [13]	$\min(\mathbf{n}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2}) > \mathbf{f}_{\mathcal{C}_1} + \mathbf{f}_{\mathcal{C}_2}$	$\mathbf{f}_{\mathcal{C}_1}$	$+\mathbf{f}_{\mathcal{C}_2}+1$	$\checkmark$	×
	PBS-cs [13]	$\mathbf{n}_{\mathcal{C}_1} > 3\mathbf{f}_{\mathcal{C}_1},  \mathbf{n}_{\mathcal{C}_2} > 3\mathbf{f}_{\mathcal{C}_2}$	max	✓	×	
	GEOBFT [12]	$\mathbf{n}_{\mathcal{C}_1} = \mathbf{n}_{\mathcal{C}_2} > 3 \max(\mathbf{f}_{\mathcal{C}_1}, \mathbf{f}_{\mathcal{C}_2})$	$\mathbf{f}_{\mathcal{C}_2} + 1^{\ddagger}$	$\Omega(\mathbf{f}_{\mathcal{C}_1}\mathbf{n}_{\mathcal{C}_2})$	×	✓
his Paper	PPCS	$\mathbf{n}_{\mathcal{C}_1} > 2\mathbf{f}_{\mathcal{C}_1},  \mathbf{n}_{\mathcal{C}_2} > 2\mathbf{f}_{\mathcal{C}_2}$	4	$(\mathbf{f}_{\mathcal{C}_1}+1)(\mathbf{f}_{\mathcal{C}_2}+1)$	×	~
	PPCS	$\mathbf{n}_{\mathcal{C}_1} > 3\mathbf{f}_{\mathcal{C}_1},  \mathbf{n}_{\mathcal{C}_2} > 3\mathbf{f}_{\mathcal{C}_2}$	$2\frac{1}{4}$	$(\mathbf{f}_{\mathcal{C}_1}+1)(\mathbf{f}_{\mathcal{C}_2}+1)$	×	~
	PLCS	$\min(\mathbf{n}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2}) > \mathbf{f}_{\mathcal{C}_1} + \mathbf{f}_{\mathcal{C}_2}$	4	$\mathbf{f}_{\mathcal{C}_1} + \mathbf{f}_{\mathcal{C}_2} + 1$	<	✓
	PLCS	$\min(\mathbf{n}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2}) > 2(\mathbf{f}_{\mathcal{C}_1} + \mathbf{f}_{\mathcal{C}_2})$	$2\frac{1}{4}$	$\mathbf{f}_{\mathcal{C}_1} + \mathbf{f}_{\mathcal{C}_2} + 1$	$\checkmark$	✓
	PLCS	$\mathbf{n}_{\mathcal{C}_1} > 3\mathbf{f}_{\mathcal{C}_1},  \mathbf{n}_{\mathcal{C}_2} > 3\mathbf{f}_{\mathcal{C}_2}$	3	$\max(\mathbf{n}_{\mathcal{C}_1},\mathbf{n}_{\mathcal{C}_2})$	<	~

# **Byzantine Cluster-Sending in Expected Constant Communication [arXiv'20]**

formalizing the problem of probabilistically sending a message from one

Byzantine cluster to another Byzantine cluster in a reliable manner,

establishing lower bounds on the complexity

of this problem under crash failures and Byzantine failures

(expected constant message complexity)

$\mathcal{C}_1$ :	R <sub>1,1</sub>	R <sub>1,2</sub>	R <sub>1,3</sub>	R <sub>1,4</sub>	R <sub>1,5</sub>	R <sub>1,6</sub>	R <sub>1,7</sub>	R <sub>1,8</sub>							
$\mathcal{C}_2$ :		¥ R <sub>2,1</sub>	¥ R <sub>2,2</sub>	¥ R2,3	₩ R <sub>2,4</sub>	$\mathbf{k}_{\mathrm{R}_{2,5}}$	¥ R <sub>2,6</sub>	R2,7							
	Ċ						$\bigtriangledown$			Protocol	Robustness	Mes (expected)	ssage Steps (worst)	0.	U.
								_		PBS-cs [13]	$\min(\mathbf{n}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2}) > \mathbf{f}_{\mathcal{C}_1} + \mathbf{f}_{\mathcal{C}_2}$	$\mathbf{f}_{\mathcal{C}_1}$	$+\mathbf{f}_{\mathcal{C}_2}+1$	✓	×
							$PBS-CS [13]  \mathbf{n}_{\mathcal{C}_1} > 3\mathbf{f}_{\mathcal{C}_1},  \mathbf{n}_{\mathcal{C}_2} > 3\mathbf{f}_{\mathcal{C}_2} \qquad \max$		$\mathbf{x}(\mathbf{n}_{\mathcal{C}_1},\mathbf{n}_{\mathcal{C}_2})$	✓	×				
										GeoBFT $[12]$	$\mathbf{n}_{\mathcal{C}_1} = \mathbf{n}_{\mathcal{C}_2} > 3 \max(\mathbf{f}_{\mathcal{C}_1}, \mathbf{f}_{\mathcal{C}_2})$	$\mathbf{f}_{\mathcal{C}_2} + 1^{\ddagger}$	$\Omega(\mathbf{f}_{\mathcal{C}_1}\mathbf{n}_{\mathcal{C}_2})$	×	~
								_	ម	PPCS	$\mathbf{n}_{\mathcal{C}_1} > 2\mathbf{f}_{\mathcal{C}_1},  \mathbf{n}_{\mathcal{C}_2} > 2\mathbf{f}_{\mathcal{C}_2}$	4	$(\mathbf{f}_{\mathcal{C}_1}+1)(\mathbf{f}_{\mathcal{C}_2}+1)$	×	~
									ape	Ppcs	$\mathbf{n}_{\mathcal{C}_1} > 3\mathbf{f}_{\mathcal{C}_1},  \mathbf{n}_{\mathcal{C}_2} > 3\mathbf{f}_{\mathcal{C}_2}$	$2\frac{1}{4}$	$(\mathbf{f}_{\mathcal{C}_1}+1)(\mathbf{f}_{\mathcal{C}_2}+1)$	×	~
								ſ	n L L L L	PLCS	$\min(\mathbf{n}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2}) > \mathbf{f}_{\mathcal{C}_1} + \mathbf{f}_{\mathcal{C}_2}$	4	$\mathbf{f}_{\mathcal{C}_1} + \mathbf{f}_{\mathcal{C}_2} + 1$	<b>~</b>	~
								[	Lhi:	PLCS	$\min(\mathbf{n}_{\mathcal{C}_1}, \mathbf{n}_{\mathcal{C}_2}) > 2(\mathbf{f}_{\mathcal{C}_1} + \mathbf{f}_{\mathcal{C}_2})$	$2\frac{1}{4}$	$\mathbf{f}_{\mathcal{C}_1} + \mathbf{f}_{\mathcal{C}_2} + 1$	✓	~
										PLCS	$\mathbf{n}_{\mathcal{C}_1} > 3\mathbf{f}_{\mathcal{C}_1},  \mathbf{n}_{\mathcal{C}_2} > 3\mathbf{f}_{\mathcal{C}_2}$	3	$\max(\mathbf{n}_{\mathcal{C}_1},\mathbf{n}_{\mathcal{C}_2})$	✓	<b>~</b>

# **Coordination-Free Byzantine Replication With Minimal Communication Costs [ICDT'20]**

formalizing the Byzantine learner problem to support efficient

analytics for blockchain applications

introducing the delayed-replication algorithm,

utilizing information dispersal techniques,

giving rise to a coordination-free, push-based, minimal communication protocol

![](_page_27_Figure_6.jpeg)

# **Coordination-Free Byzantine Replication With** Minimal Communication Costs [ICDT'20]

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introducing the delayed-replication algorithm,

utilizing information dispersal techniques,

giving rise to a coordination-free, push-based, minimal communication protocol

![](_page_28_Figure_6.jpeg)

## **Permissioned Blockchain Through the Looking Glass:** Architectural and Implementation Lessons Learned [ICDCS'20]

Single-threaded Monolithic Design Out-of-ordering Consensus Communication De-coupled Ordering and Execution Off-Chain Memory Management Expensive Cryptographic Practices (DS vs. MAC) Smart Contracts Code Generation (Pre-compilation)

![](_page_29_Figure_2.jpeg)

# **Revisit Resiliency** (Graduate Student Experiment Continues)

Mount Tallac, Lake Tahoe 12.1 Miles Long 3,931 Feet Elevation Gain (9,738 Feet at Summit)

![](_page_31_Picture_1.jpeg)

![](_page_31_Picture_2.jpeg)

# **Fostering Resiliency**

#### (Offering Stress Management and Well-Being Courses at UC Davis)

![](_page_32_Picture_2.jpeg)

![](_page_32_Picture_3.jpeg)

Spring 2020

Tamarkoz®.

ECS 298 (CRN 66553):

**Days: Wednesdays** 

**INSTRUCTORS:** 

Mohammad Sadoghi, Ph.D. Nasim Bahadorani, DrPH.

Time: 7:00 pm - 8:00 pm

Graduate Survival Kit

Learn the foundation & working

knowledge of stress reduction based on a unique heart-centered

meditation practice referred to as

The M.T.O. Tamarkoz® method is

the art of self-knowledge through

concentration and meditation.

The California Aggie, April 6, 2020

BE BALANCED

![](_page_33_Picture_0.jpeg)

![](_page_33_Picture_1.jpeg)

# THANK YOU

## For Complete References

![](_page_33_Picture_4.jpeg)

![](_page_33_Picture_5.jpeg)